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Flow based intrusion detection system

*A Research Submitted in Partial fulfillment for the Requirements of
the Degree of B.Sc. (Honors) in Electronics Engineering*

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الإستهلال

بِسْمِ اللَّهِ الرَّحْمَنِ الرَّحِيمِ

قال تعالى:

﴿وَلَقَدْ كَرَّمْنَا بَنِي آدَمَ وَحَمَلْنَاهُمْ فِي الْبَرِّ وَالْبَحْرِ وَرَزَقْنَاهُمْ مِّنَ الطَّيِّبَاتِ
وَفَضَّلْنَاهُمْ عَلَى كَثِيرٍ مِّمَّنْ خَلَقْنَا تَفْضِيلًا﴾

صَدَقَ اللَّهُ الْعَظِيمَ

سورة الإسراء الآية 70 ﴿﴾

DEDICATION

***To our parents, who directed my first step in the right way, taught us to trust in Allah, believe in our self and that so much could be done with little...*

***To our brothers and sisters...*

*** To our friends who are always ready to help...*

***To our colleges...*

We dedicate this work....

ACKNOWLEDGMENTS

Our grateful thanks firstly to our God who guided us to the strait way in our life. Then many thanks and appreciations are extended to our Supervisor **Dr.Ahmed AbdAlla**for his valuable advices and endless efforts to make this work come into reality. A lot of thanks is given to the School of Electronic Engineering.

ABSTRACT

The use of packet based NIDS is expensive because each packet must be inspected deeply. This research provides solution for discovering network attacks in efficient manner using flow based network intrusion detection system. The designed system closely monitors the internet traffic based on some time-based aggregated traffic (TAT) features to determine existence of brute-force attack .These TAT features are extracted from a previously dataset of NetFlow records using a C code program. The designed system provides a property of discovering attacks with undefined signature (unknown attacks).The obtained result shows reduction in false alarm and high level security provided by this system.

المستخلص

استخدام نظام اكتشاف الاختراقات استناداً علي محتويات الحزمه مكلف لأنه يتطلب فحص دقيق لكل محتويات الحزمه. هذا البحث يوفر حل لاكتشاف الاختراقات التي تحدث في الشبكات بطريقة ذات كفاءة وذلك باستخدام نظام اكتشاف الاختراقات علي الشبكة استناداً على بعض المعاملات. هذا النظام يقوم بصوره محكمة بمراقبة حركة البيانات على الانترنت بناءً على استخلاص خصائص لبيانات مجمعة في فترات زمنية معينة عن طريق برنامج سي باستخدام مجموعة بيانات من سجلات بروتوكول محدد مسبقاً لتحديد ما اذا كان هنالك اختراق ام لا . النظام بتوفير خاصية اكتشاف الاختراقات التي لا يكون لديها نمط محدد و معروف ، النتائج التي تم الحصول عليها تشير الى انه تم تقليل نسبة الانذار الخاطئ وتوفير درجة عالية من الأمان باستخدام هذا النظام.

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LIST OF ABBREVIATION

IDS	Intrusion Detection System
NIDS	Network Intrusion Detection System
HIDS	Host Intrusion Detection System
IPFIX	Internet Protocol Flow Information Export
DOS	Denial Of Service
DDOS	Distributed Denial Of Service
FPR	Fouls Positive Rate
TPR	True Positive Rate
AVG	Average
STDEV	Standard Deviation

Chapter one

INTRODUCTION

1.1 Back ground

Intrusion Detection can be defined as “ security system that monitors computer systems and network traffic and analyzes that traffic for possible hostile attacks originating from outside the organization and also for system misuse or attacks originating from inside the organization”[1]. It detects unwanted exploitation to computer system, both through the Internet and Intranet.

In general, we can divide IDSs into two basic classes based on their position in the network: host-based IDSs (HIDSs) and network-based (NIDSs). IDSs also can be classified based on its detection model into two categories: signature-based and anomaly-based. Signature-based IDS maintains a database of known intrusion technique (attack signature) and detects intrusion by comparing behavior against the database. Disadvantage of this technique is ineffective against previously unseen attacks and hence it cannot detect new and unknown intrusion methods as no signatures are available for such attacks. Anomaly based NIDS monitors network traffic and compares it against an established baseline of normal traffic profile. The baseline characterizes what is "normal" for the network - such as the normal bandwidth usage, the common protocols used, correct combinations of ports numbers and devices to detect malicious traffic, parameters like number of connection request, number of rejected connection ,average packet size, flag values present in the packet headers are used. But due to the dynamic nature of network traffic and application access, this type of detection can generate false alarms.

Required challenge in anomaly detection is the volume of data for analysis. Collection and analyses of network traffic information at packet level for a high-speed network to provide accurate result in real-time is a

difficult task. In case of high speed network, packet-level traffic monitoring is expensive because of deep packet inspection requirement and those intrusion detection systems can detect only known attacks based on signatures. Due to these reasons, flow based traffic monitoring and Anomaly detection is important.

1.2 Problem statement

Collection and analyses of network traffic information at packet level for a high-speed network to provide accurate result in real-time is a difficult task. In case of high speed network, packet-level traffic monitoring is expensive and it is very time consuming because of deep packet inspection requirement.

Signature-based IDSs cannot detect unknown attacks, either because the database is out of date or because no signature is available yet, and they cannot detect some type of attacks such as scan, flood, DoS and DDoS because they have no signature pattern and they appear as abnormal behavior in the network.

1.3 Proposed solution

The proposed solution for the problems is by using a flow based anomaly network intrusion detection system. The system use the "flow-level NIDSs", in which rather than looking at all packets going through a network link, it looks at aggregated information of related packets of network traffic in the form of flow, so the amount of data to be analyzed is reduced.

Also the system use the "Anomaly-based or behavior-based method" which works by building a model of normal traffic data pattern during a training phase, then it compares new inputs to the model. A significant

deviation (change) is marked as an anomaly (abnormal or intrusion), so this method is able to detect unknown attacks.

1.4 Research aim

The aim of this project is to design and implement a flow based anomaly network intrusion detection system that can detect Denial of Service, flooding, host and port scan attacks.

1.5 Research Outline

This thesis is organized as follow: Chapter two introduces the network intrusion detection system and its types. In addition, the chapter reviews types of attacks. In chapter three, the design of "flow based network intrusion detection system" was demonstrated. In chapter four the results were showed and discussed. In chapter five, the conclusion and recommendations were presented.

Chapter two

Literature Review

2.1 Introduction

The rapid proliferation of computer networks has changed the prospect of network security. An easy accessibility condition cause computer network's vulnerable against several threats from hackers. Threats to networks are numerous and potentially devastating. Recent reports on Internet security breaches indicate that the frequency and he damage costs are continuously rising. Up to the moment, researchers have developed Intrusion Detection Systems (IDS) capable of detecting attacks in several available environments.

2.2 Intrusion detection system

An IDS is best defined as software or hardware used to detect unauthorized traffic or activities that are against the allowed policy of a given network [2]. IDS do it by collecting data from network and analysis of transmitted packets inside the network. But generally IDSs do not act operative reaction against occurred attacks. IDS have many classification based on several aspects.

2.2.1 Based on detection model

If a system bases the detection on a definition of *normal* behavior of the target system, it is called *behavior-based*. If it matches the input data against a definition of an attack, it is known as *knowledge-based*. In literature, the community usually refers to these classes with the names of *anomaly-based* and *misuse-based* [3].

2.2.1.1 Signature based

The signature-based IDSs, also named “misused-based”, works similar to anti-virus software. It employs a signature (pattern that correspond to a known threat) database of known attacks, and if a

successful match with current input, an alert is raised. A well-known example of this type is Snort [4] which is an open source IDS that monitors network by matching each packet it observes against a set of rules.

Signature based method is very useful for known attacks, Although Signature-based IDSs cannot detect unknown attacks, either because the database is out of date or because no signature is available yet, it has low false alarms (high accuracy), Signature based systems are reactive, in that they combat against known attacks, that have already affected and damaged a number of systems before being identified.

2.2.1.2 Anomaly based

Anomaly detection is an active area in network intrusion detection research which was originally proposed by Denning [5]. It can detect various types of intrusion based on the deviation in the normal usage of network and this has an advantage over signature based technique.

Anomaly-based or behavior-based IDS works by building a model of normal traffic data pattern, then it compares new inputs to the model. A significant deviation (change) is marked as an anomaly (abnormal or intrusion). Anomaly-based is able to detect unknown attacks but it suffers from producing false alarms [4]. It refers to finding out the abnormal pattern of traffic or abnormal behavior from network or system.

2.2.2 Based on their position in the network

Most traditional intrusion detection systems (IDS) take either a network- or a host-based approach to recognizing and deflecting attacks. Each approach has its strengths and weaknesses.

2.2.2.1 Host based

Host-based intrusion detection systems are aimed at collecting information about activity on a particular single system, or host[6] . The term “host” refers to an individual computer.

2.2.2.2 Network based

Network-based intrusion detection systems offer a different approach; “These systems collect information from the network itself,” [6] rather than from each separate host. information is collected from the network traffic stream, as data travels on the network segment [6].

Network-based systems are extremely portable. They only monitor traffic over a specific network segment, and are independent of the operating systems that they are installed on. There are two methods basis on the source of data to be analyzed in NIDSs:

- **Packet based**

In packet-based, also named “Deep Packet Inspection” (DPI), the combination of header and payload scan determines whether a packet is an intrusion or not. Incoming packets are scanned and every single rule of the database is checked against it as shown in figure 2-1. The database rules include thousands of signatures and patterns of attacks.

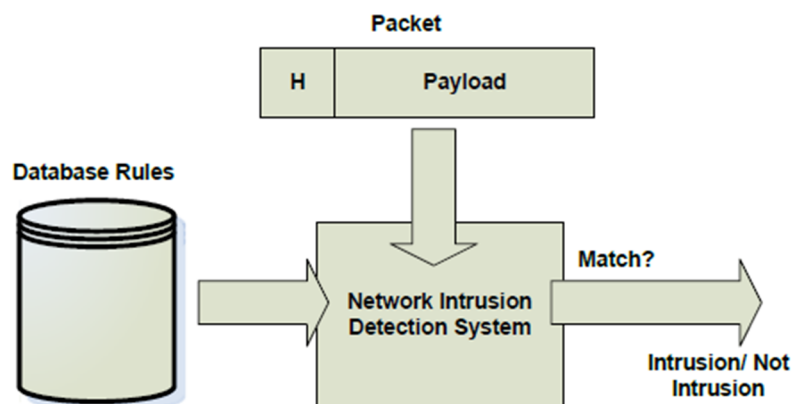


Figure 2.1: packet based IDS

The main advantage of packet-based approach is that all common kinds of known attacks and intrusions practically can be detected if the data source deliver entire network packet for analysis. However, it cannot detect unknown attack since it compare with predefined and known malicious signatures [4].

However, systems that are capable of monitoring every packet on a high-speed network are very expensive and high resource consumption. Moreover, a drop of packets will occur if the NIDSs speed is not high enough to let the analysis process be done.

In signature-base detection, as the number of attacks increase, the number of malicious (intrusion) signatures increase in the database in NIDSs. Usually, these databases contain hundreds or thousands of signatures. NIDS has to add these new signatures into its signature list quickly without disturbing its main function of detecting intrusion. NIDSs then search for these signatures in network traffic to detect intrusions. To detect signatures, all network traffic must be compared against every signature to identify if a match exists or not. Therefore, the efficiency in accessing such database for analyses is also critical. Another issue is that signature matching is impossible for most cases of encrypted payload, degrading the detection performance of NIDSs. A comprehensive evaluation of packet-based performance with high volume NIDSs is presented in [4].

- **Flow based**

For high speed networks, it is important to explore alternative to packet-based inspection for efficient NIDSs. One option that currently attracts the attention of researchers is flow-based intrusion detection. Flow-based technique is widely deployed as data source in applications

like network monitoring, traffic analysis and security [4]. This method is characterized by flow data or network flow. Flows don't provide any packet payload unlike packet-based approach, as shown in figure 2.2.

It rather relies on information and statistics of network flows. A flow can be defined as a unidirectional data stream between two computer systems where all transmitted packets of this stream share the following characteristics: IP source and destination address, source and destination port number and protocol value [4]. Nowadays special measurement systems are able to provide other characteristics in addition to the above, for instance:

- The number of packets and amount of bytes transferred in a flow.
- The start and end time of a flow (in milli-second).

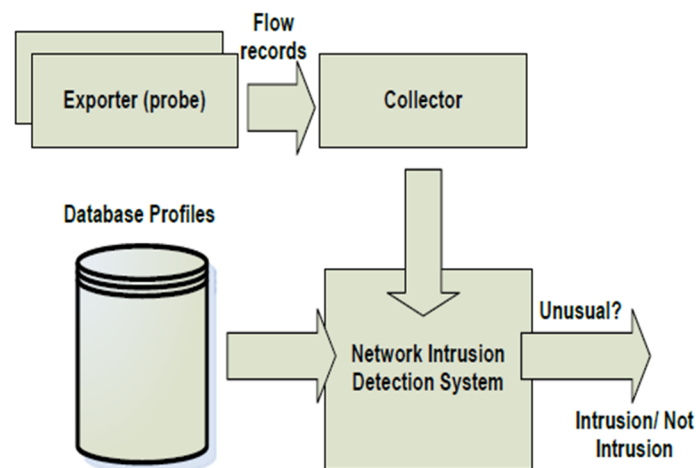


Figure 2.2: flow based IDS

2.3 network attacks

Network and computer attacks have become pervasive in today's world. Any computer connected to the Internet is under threat from viruses, worms and attacks from hackers. Home users, as well as business

users, are attacked on a regular basis. Thus the need to combat computer and network attacks is becoming increasingly important.

Flow-based intrusion detection, since it relies only on header information, can address only a subset of the attacks presented above. In particular, the research community currently provides approaches to detect the following classes of attacks:

- Denial of Service.
- Scans.
- flooding.

2.3.1.1 Denial of Service attack

A Denial of Service attack attempts to slow down or completely shut down a target so as to disrupt the service and deny the legitimate and authorized users can access [7]. Such attacks are very common in the Internet where a collection of hosts are often used to bombard web servers with dummy requests illustrated in Figure 2.3. Such attacks can cause significant economic damage to ecommerce businesses by denying the customers an access to the business. There are a number of different kinds of DoS attacks, some of which are mentioned below. It is shows a denial of service attack (DDoS in this case), wherein an attacker uses a number of compromised hosts to attack a given victim.

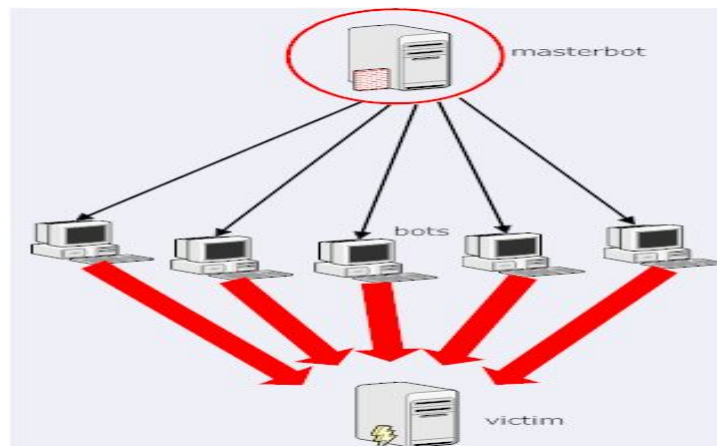


Figure 2.3: Denial of Service attack

- **Flaw Exploitation DoS Attacks**

In such attacks, an attacker exploits a flaw in the server software to either slow it down or exhaust it of certain resources. Ping of death attack is one such well known attack. A ping of death (POD) is a type of attack on a computer that involves sending a malformed or otherwise malicious ping to a computer. A ping is normally 64 bytes in size (or 84 bytes when IP Header is considered); many computer systems cannot handle a ping larger than the maximum IP packet size, which is 65,535 bytes. Sending a ping of this size can crash the target computer. Some limitations of the protocol implementation also lead to vulnerability which can be exploited to implement DoS attacks such as DNS amplification attack, which uses ICMP echo messages to bombard a target. For these attacks, a signature can be devised easily, such as to determine a ping of death attack a NIDS needs to check the ping flag and packet size.

- **Flooding DoS Attacks**

In a flooding attack, an attacker simply sends more requests to a target than it can handle. Such attacks can either exhaust the processing capability of the target or exhaust the network bandwidth of the target,

either way leading to a denial of service to other users. DoS attacks are extremely difficult to combat, as these do not exploit any vulnerability in the system, and even an otherwise secure system can be targeted. A more dangerous version of DoS attack is called Distributed Denial of Service attack (DDoS), which uses a large pool of hosts to target a given victim host. A hacker (called botmaster) can initiate a DDoS attack by exploiting vulnerability in some computer system, thereby taking control of it and making this the DDoS master (Figure 2-3). Afterwards the intruder uses this master to communicate with the other systems (called bots) that can be compromised. Once a significant number of hosts are compromised, with a single command, the intruder can instruct them to launch a variety of flood attacks against a specified target [7].

2.3.1.2 Scanning Attack

In such attacks, an attacker sends various kinds of packets to probe a system or network for vulnerability that can be exploited. When probe packets are sent the target system responds; the responses are analyzed to determine the characteristics of the target system and if there are vulnerabilities (illustrated in Figure 2-4) where a single attack host scans a number of victims. Thus scanning attack essentially identifies a potential fused which yields these information:

- The network topology.
- The type of firewall used by the system.
- The identification of hosts that are responding.
- The software, operating systems and server applications that are currently running.
- Vulnerabilities in the system.

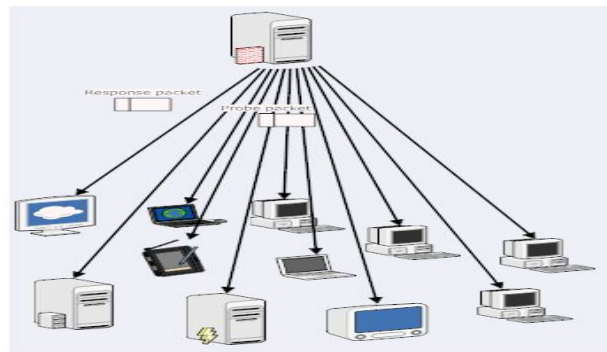


Figure 2.4: scan attack

Once the victim is identified, the attacker can penetrate them in a specific way. Scanning is typically considered a legal activity and there are a number of examples and applications that employ scanning. The most well known scanning applications are Web search engines. On the other hand independent individual ay scan a network or the entire Internet looking for certain information, such as a music or video file. Some well-known malicious scanning include Vertical and Horizontal port scanning, ICMP (ping) scanning, very slow scan, scanning from multiple ports and scanning of multiple IP addresses and ports. NIDS signatures can be devised to identify such malicious scanning activity from a legitimate scanning activity with fairly high degree of accuracy [7].

2.3.1.3 Flooding attack

- **Ping flood and ping of death:**

Ping flood is similar to Smurf where in the victim is bombarded with thousands of ping packets. In Ping of death, the victim is sent corrupt packets that could crash the system [8]. Smurf and ping floods are very easy to craft and any novice attacker could do it with ease. These attacks could cause considerable damage in small Local Area Networks.

- **UDP flood:**

UDP flooding is similar to ping flood. Here instead of ping packets, UDP packets are bombarded against the server. UDP could be a lot more effective than ICMP in smaller networks as the size of the UDP packets are enormous. The packet size could be set up to 65000 bytes which could easily flood a given Ethernet network when multiple zombies are set up. This project has analyzed all the above described attacks and has brought down some interesting observations.

2.4 Related work

Recently, instead of packet based analysis, flow based security analysis and anomaly detection are getting attention from many researchers. Mayung et al [4] suggests that by aggregating packets of the identical flow, one can identify the abnormal traffic pattern that appears during attack. They formalize detection function for attack detection, which are composed of several traffic parameters and constant value.

Another work based on flow monitoring is explained in[9] which work on monitoring the four predefined metrics that capture the flow statistic of the network. This method is capable to detect UDP flood, ICMP flood and scanning, by using Holt-Winters Forecasting technique. This technique makes projection about future performance based on historical and current data of the network. The prediction which comes out by this technique may arise false alarms because network behavior is not static.

Anomaly-based detection stated by [10] analyses user behavior and the statistics of a process in normal situation, and it checks whether the system is being used in a different manner. In addition [10] has described that this technique can overcome misuse detection problem by focusing on normal system behavior rather than attack behavior. However [10]

assume that attacks will result in behavior different from that normally observed in a system and an attack can be detected by comparing the current behavior with pre-established normal behavior. This detection approach is characterized by two phases which is the training phase and detection phase. In training phase, the behavior of the system is observed in the absence of attack, and machine learning technique is used to create a profile of such normal behavior. In detection phase, this profile is compared against the current behavior, and deviations are flagged as potential attacks. The effectiveness of this technique is affected by what aspect or a feature of system behavior is learnt and the hardest challenge is to be able to select the appropriate set of features. The advantage of this detection technique is that it can detect new intrusion method and capable to detect novel attacks. However, the disadvantage is that it needs to update the data (profiles) describing the user's behavior and the statistics in normal usage and therefore it tend to be large and therefore need more resources, like CPU time, memory and disk space. Moreover, the malware detector system often exhibit legitimate but previously unseen behavior, which leads to high rate of false alarm.

In the current trend, few researches such as [10]and[7] have been found to manipulate this detection technique by combining either Signature-based with Anomaly-based detection technique(Hybrid-SA) in order to develop an effective malware detector's tool.

A hybrid-based IDS stated by [4] make a tradeoff between availability of limited data of flow-based techniques, which have negative effect on accuracy of NIDSs, and full data of packet-based which lead to a higher resources consumption. Therefore flow-based detection should not substitute the packet inspection one. However, combination approach to combine both approaches to power their advantages and overcome

their drawbacks is proposed. Flow-based technique takes advantage of packet-based technique to reduce false alarms. We therefore expect a potential in mixing both approaches to detect at least the same quantity of attacks while consuming less resources.

2.5 Summary

There are problem that some attacks have traffic patterns that cannot be characterized by only one flow. To detect this type of attack, we need traffic information that can identify traffic patterns. Aggregating related flows can generate this information, which is called traffic pattern data. By examining parameters of traffic pattern data we can discover traffic used in attacks, such as flooding and scanning.

Chapter Three

Research Methodology

3.1 Introduction

This chapter describes the implementation details of developing and implementation of the proposed flow based anomaly detection system.

3.2 Research Activities

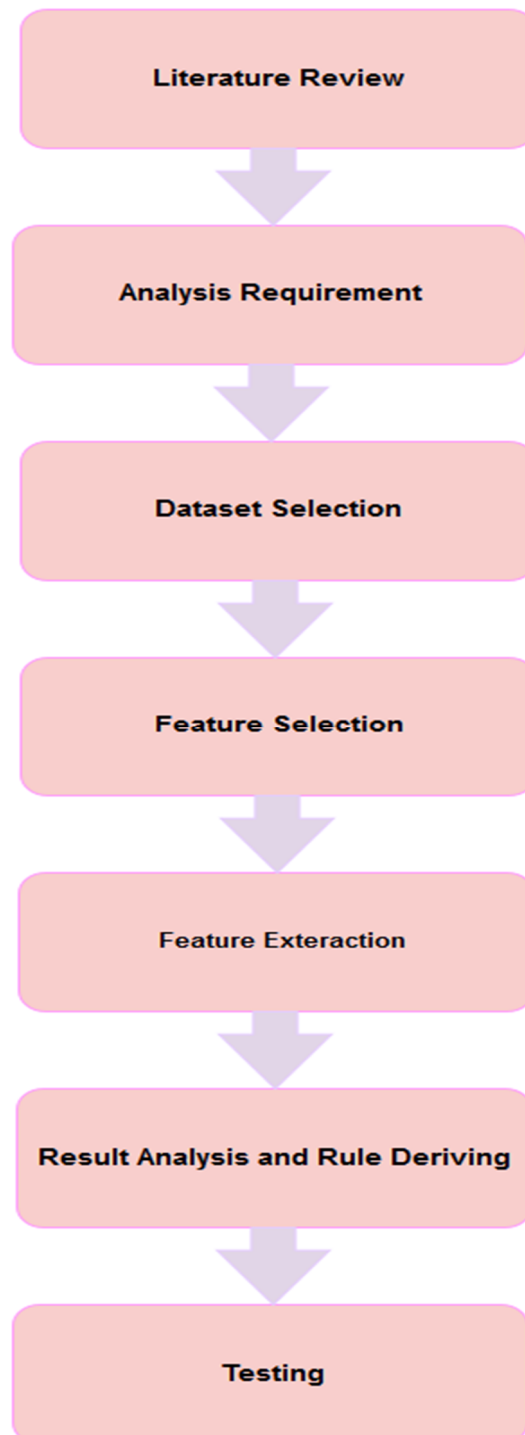


Figure 3.1: Research activities sequence

- **Literature Review**

We reviewed a number of scientific papers related with our proposed system and follow major related works, Theories and hypotheses mentioned in the literature.

- **Analyze requirements**

Specification of Dataset was been determined.

- **Data selection**

The selection of dataset depends on [11].

- **Feature selection**

We select some feature from Dataset.

- **Feature extraction**

We extract the selected feature in previous step for TAT instead of flow record.

- **Result analysis and rule deriving**

Analyze results for labeling TAT to obtain good result.

- **Testing**

Training rules with labeling TAT record to improve accuracy of rules.

3.3 Flow based anomaly detection system

This section is divided into two parts .The first part describes the characteristics of our selected dataset as prerequisites for system ,The second part describes the implementation details of designing an effective NIDS that is capable to detect brute-force attacks. In the following subsections we will describe and discuss all processes and implementation details.

3.3.1 Dataset

Our project depend on pre-prepared dataset by[11], [11] explained clearly and in details all the processes and procedures taken to create flow-level dataset. This dataset has ten attributes, but we select eight of them, as illustrated in Table 3.1.

Table 3.1: Final Eight Attributes selected for datasets

	attribute	attribute Description
1	Stime	Unix time of the first packet in the flow
2	etime	Unix time of the last packet in the flow
3	SrcIP	source IP
4	DstIP	destination IP
5	SrcPT	TCP/UDP source port number or equivalent
6	DstPT	TCP/UDP destination port number or equivalent
7	pkt	Number of Packets in the flow
8	Pyt	Number of Layer 3 bytes in the packets of the flow

Labeling means determines if the flow record is malicious or free attack. As [11] mention, each malicious flow is assigned three digits to identify the specific attack type. Table 3.2 shows this labeling method. A label of “223” means that the attack is a flood because the first digit is “2”, DDoS because the second digit is “2” and the used protocol is UDP because the third digit is “3”.

Table 3.2: Three Digit labeling method of malicious flow

FIRST DIGIT: main class attack			"1" ↓ scan			"2" ↓ flood	
SECOND DIGIT specific type into main class			"1" ↓ network scan	"2" ↓ port scan	"3" ↓ protocol scan	"1" ↓ DoS flood	"2" ↓ DDoS flood
THIRD DIGIT used protocol	1→	ICMP	☑ 111			☑ 211	☑ 221
	2→	TCP	☑ 112	☑ 122		☑ 212	☑ 222
	3→	UDP	☑ 113	☑ 123		☑ 213	☑ 223
	0→	OTHER	☑ 110	☑ 120	☑ 130	☑ 210	☑ 220

3.3.2 Flow-Level Brute-Force Attack Detector

The objective of this part is to aggregate previously prepared NetFlow records in fixed time windows and extract traffic features that give optimum detection result according to the selected detection method and finally to come out with rules for detection and identification of brute-Force Attack in near real time. The proposed system is presented in Figure 3.2. As illustrated, the system has the following main sequential processes Time-based Aggregation of Traffic (TAT), labeling TAT records, developing of a detection model based on the TAT records and testing.

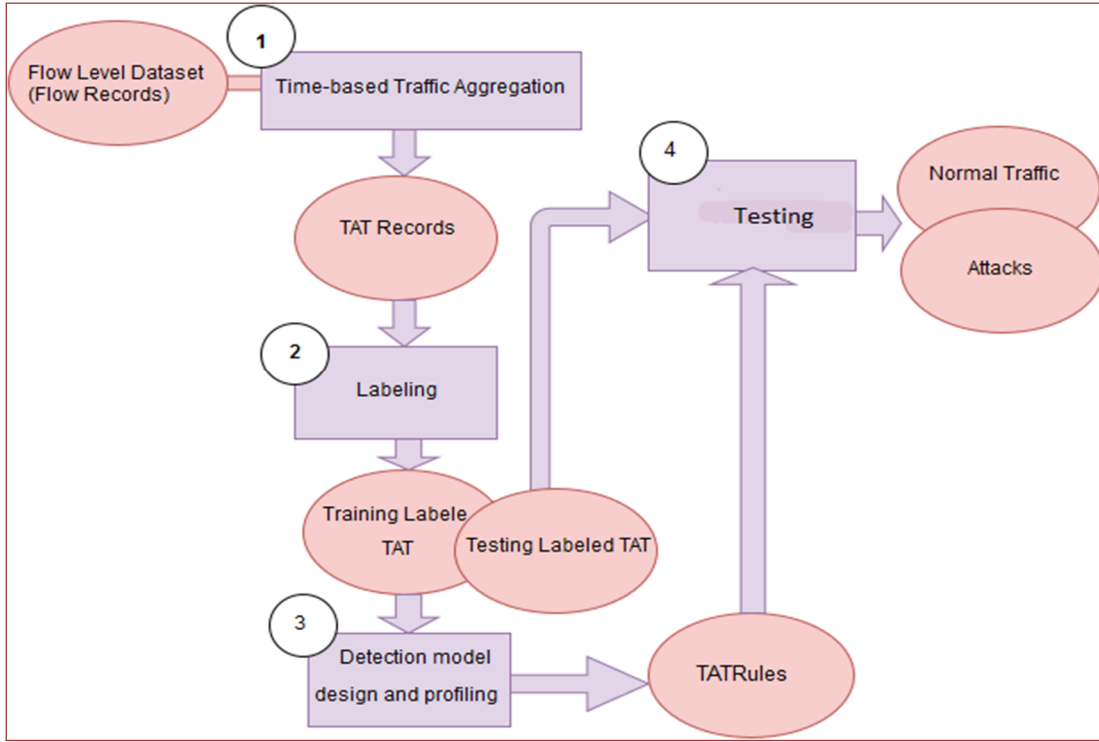


Figure 3.2: Designing and Testing of the Flow-Level Brute-Force Detector.

3.3.2.1 Time-based Aggregation of Traffic

The *flow-level dataset* entity which appears in Figure 3.1 as input of the *Time-based Aggregation of Traffic* process is the same brute-force dataset which had been described in previous section. In this process, we group NetFlow records of the brute-force dataset into window of t second then we extract various statistical metrics which we call Time-based Aggregated Traffic (TAT) features. Figure 3.3 represents a conceptual diagram showing this process. We chose time window for our experiments 60 sec, we wrote a C program to implement the described process (see **Appendix A**). Figure 3.4 illustrates the flow chart of that program.

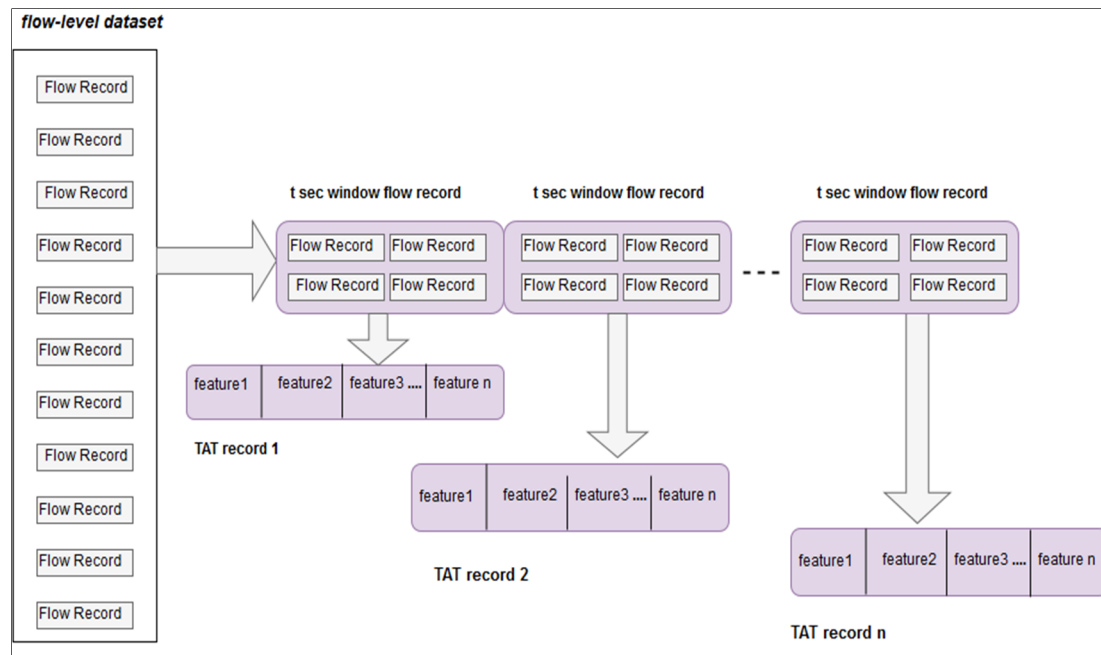


Figure 3.3 Time-based Aggregation of Traffic

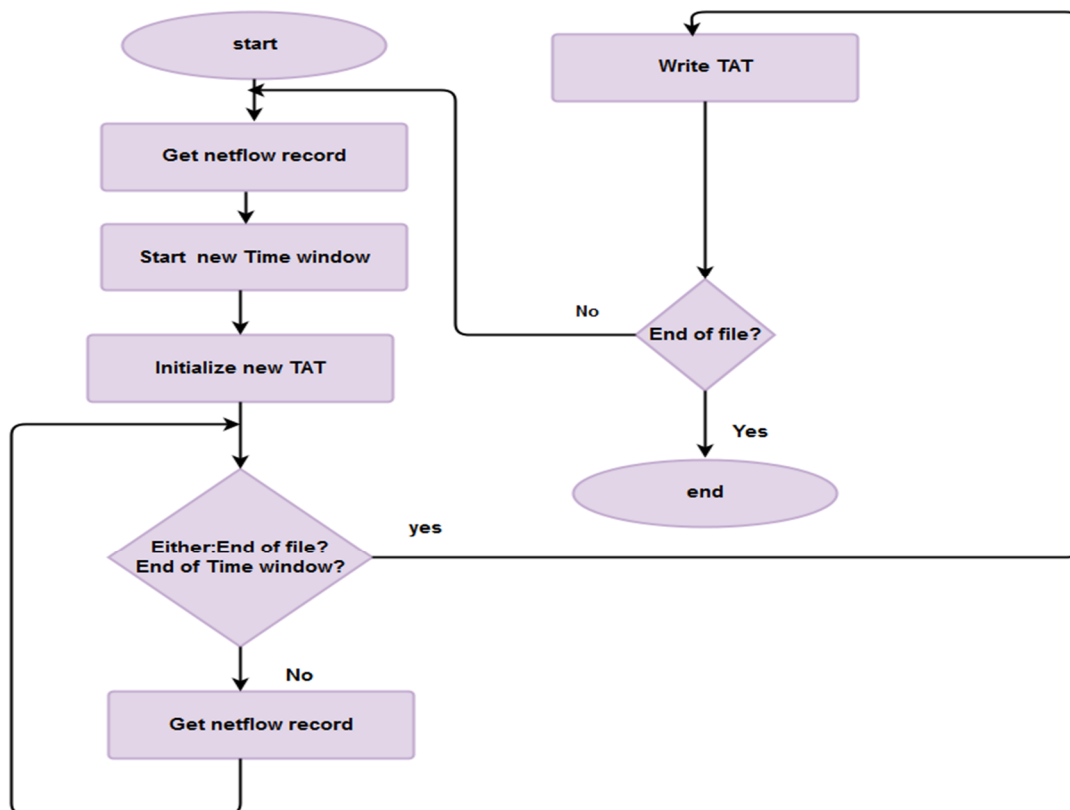


Figure 3.4: flow chart of Time Aggregation of Traffic

3.3.2.2 Labeling TAT Records

In labeling TAT there is a problem of how to classify a TAT record when it contains both classes of benign and malicious flows. To deal with this problem, we chose to label a whole TAT record as malicious if the record satisfies that the number of malicious flows is 50 flows or greater. We modified the previous written program that implements Traffic Aggregation labeling for each TAT (see **Appendix A**). The flow chart of the modified program is illustrated in Figure 3.5

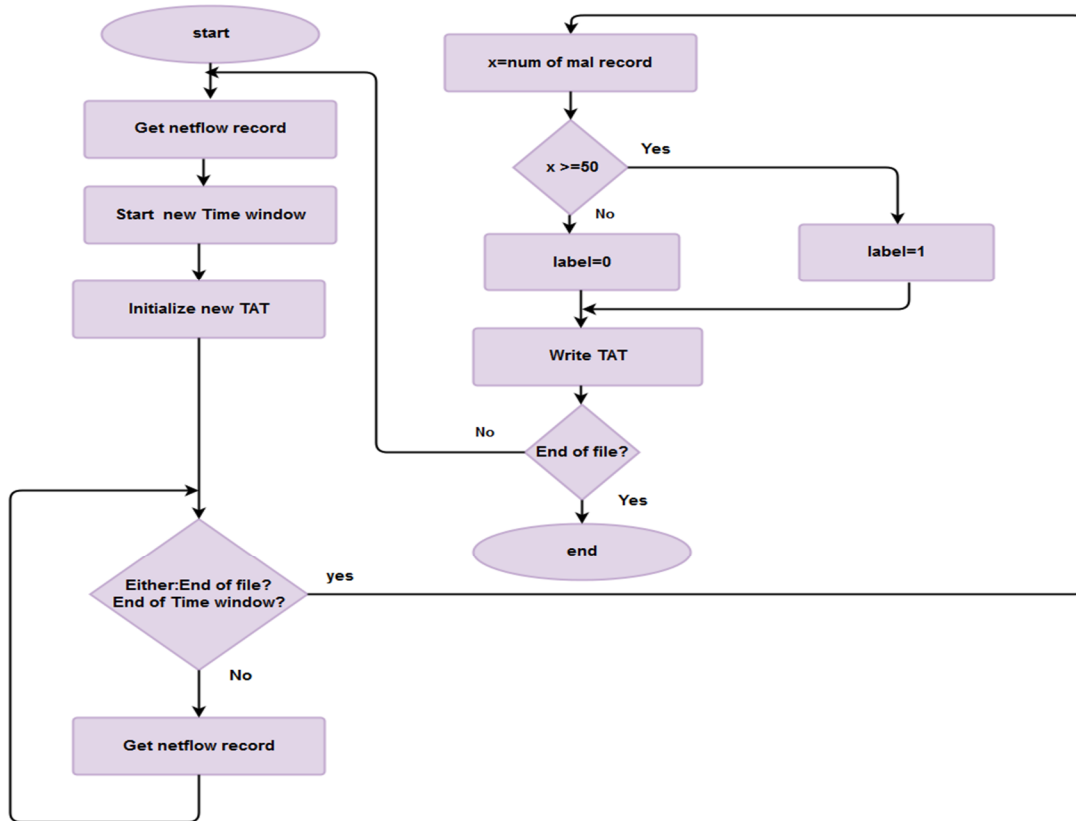


Figure3.5: Flow Chart of Traffic Aggregation of Traffic and TAT record Labeling.

3.3.2.3 Detection Model Design and Profiling

Design of the detection model is done in three sequential steps. The first phase is the selection of significant TAT attributes that reflect traffic variations of brute-force attacks. The second step is determining threshold values that differentiate normal and malicious TAT records. The final step is to formulize detection rules from combinations of these threshold values.

- **Selection of Significant TAT attributes**

We may assess the effect of each feature in detecting attacks by analyze the labeled TAT records to study the effect of each attack over TAT attributes.

- **TAT statistics and Profiling**

For detection process we calculated statistical features (AVG, STDEV) for each TAT attribute to characterize, differentiate and identify normal TAT records and malicious TAT records.

- **Threshold Values and Detection Rules**

We use the combination of (AVG, STDEV) for TAT attribute as threshold, then detection rule will explain in next chapter.

3.3.2.4 Testing

This testing is done on labeled TAT records to check the detection method and identification accuracy in terms of true positive and false positive for brute force attack. We implement the constructed attack detection rules showed in the previous section. We checks a TAT record whether it is benign or malicious.

Chapter four

Results

4.1 Introduction:

In this chapter we illustrate result and output of each phase in proposed system.

4.2 Result of Time-based Aggregation Traffic

Table 4.1 lists and describes the full attribute set (TAT record fields) extracted out of one time window of NetFlow records.

Table 4.1: Extracted Features of Time Window NetFlow Records

	TAT feature	Description
1	#flows	number of flows
2	#2pflow	number of flows that contain 1 to 2 packets
3	#pkts	number of packets
4	#byts	number of bytes
5	#srcIP	number of distinct source IP addresses
6	#dstIP	number of distinct destination IP addresses
7	#srcPt	number of distinct source Ports
8	#dstPt	number of distinct destination Ports

4.3 Result of Labeling TAT Records

After labeling phase the TAT records statistics shown in Table 4.2 calculated from result obtained from C program (see Appendix B).

Table 4.2: Statistics of 60 seconds window TAT records

TAT records statistics	60 sec TAT
Malicious TAT records	49
free attack TAT records	74

4.4 Result of Detection Model Design and Profiling

Statistical from Significant TAT attributes and threshold used in rule for detection illustrated in Table 4.3.

Table 4.3: TAT attributes threshold

Attribute	Threshold
# flows	18417
# srcIP	1408
# srcPt	7534
# pkt2flow	41

Where:

$$\text{pkt2flow} = 2\text{pflow} / \text{flows}.$$

After experiment and observation, we notice that the features in Table 4.3 have more effect as shown in figure 4.1 this features had been used to formed suitable rule.

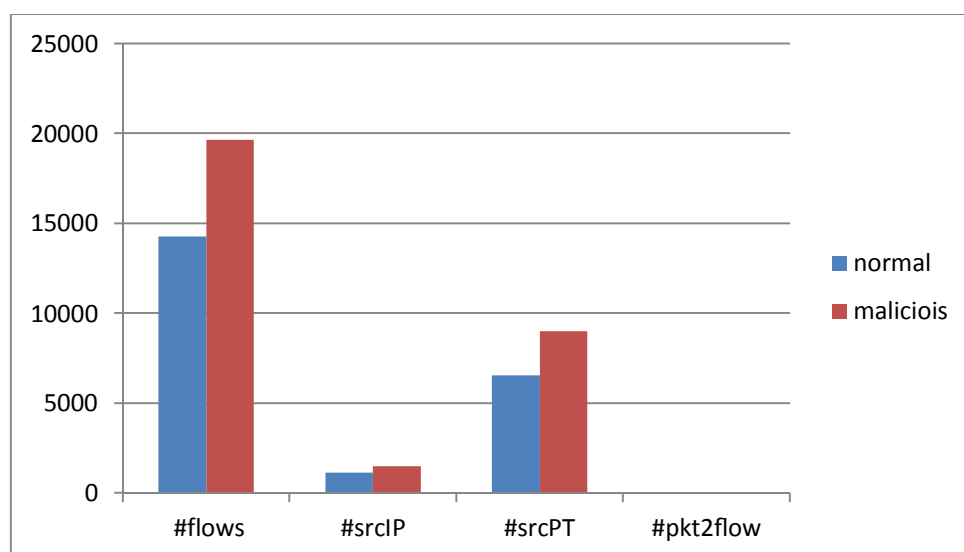


Figure 4.1: TAT Attributes for Normal and Brute force attack (Average Value)

The Rule:

IF((flows>=18417 AND srcIP>=1408) OR (pkt2flow <=41 AND srcPt >=7534))

Label =1

else

Label=0

4.5 Result after Testing Rule

After applied rules of detection method with labeled TAT record we obtain the result shown in Table 4.4

Table 4.4: accuracy of detection rule

TPR	FPR
80%	8%

Chapter Five

Conclusion

5.1 Conclusion

A flow based detection system is proposed in this research. The proposed system provides a generic solution for detecting network anomalies like scan and flood (Brute force attack) for high speed network. The achieved results were satisfactory with ratio 86.99%.

5.2 Recommendations

This research has provided report on the design process of our system; definitely we need more study to improve reliability, capability and accuracy of our system. The following improvements can be recommended for possible future work:

- Another statistical method can be used, like: chi square .
- More features can be used, or combination of any effective features.
- This system can be implemented online.
- Determining the classification of the attacks types can be added.
- Using more time windows and select the optimum one, will give more accurate result.

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APPENDIX

A.APPENDIX A

CODE OF AGGREGATION AND LABELING

```
#include<stdio.h>
#include<stdlib.h>

#define SIZE 100000 //maximum size of array
FILE *rf, *wf,*wr;
unsigned long no_flow
,sum_pkts,sum_bytes,pkt2,s[SIZE],y[SIZE],x[SIZE],
v[SIZE],usrcIP,udstIP,usrcPT,udstPT,
foundIPadrs,foundIPdst,foundptdrs,foundptadst,e,o
o;

int j=0,n,a,b,c,r,label;

struct FlowRecord {
    unsigned long    flowId;
    unsigned long    stime;
    unsigned long    stmsec;
    unsigned long    etime;
    unsigned long    etmsec;
    unsigned char    protocol;
    unsigned long    srcIP;
    unsigned long    srcPT;
    unsigned long    dstIP;
```

```

        unsigned long    dstPT;
        unsigned long    pkts;
        unsigned long    bytes;
        unsigned long    clss;
    };
    struct FlowRecord FR;

    void R_file();
    void W_file();
    void parameter();
    void writeTAT();
    void initializeTAT();

    int main(void)
    {
        if ((rf=fopen("D:/eflows2014-11-
13_15_43_05clss.csv", "r"))==NULL)//reading file
        {
            printf( "Read file could not be opened\n" );
            exit(1);
        }

        if ((wf=fopen("D:/writingfile.txt",
"w"))==NULL)//writing file
        {
            printf( "Write file could not be opened\n" );
            exit(1);
        }
    }

```

```

R_file();
while(!feof(rf)){
    initializeTAT();
    do{

        W_file();//calling function
        parameter();//calling function
        R_file();//calling function

    } while( !feof(rf) &&(r +60 >
FR.etime)); //feof was test

writeTAT();
}

fclose(rf);
fclose(wf);

return 0;
}

void R_file()//function read from file and write
on other file
{
fscanf(rf,"%lu",&FR.flowId);//reading
fscanf(rf,"%lu",&FR.stime);
fscanf(rf,"%lu",&FR.stmsec);
fscanf(rf,"%lu",&FR.etime);

```

```

fscanf(rf,"%lu",&FR.etmsec);
fscanf(rf,"%u",&FR.protocol);
fscanf(rf,"%lu",&FR.srcIP);
fscanf(rf,"%lu",&FR.srcPT);
fscanf(rf,"%lu",&FR.dstIP);
fscanf(rf,"%lu",&FR.dstPT);
fscanf(rf,"%lu",&FR.pkts);
fscanf(rf,"%lu",&FR.bytes);
fscanf(rf,"%lu",&FR.class);
}

```

```

void initializeTAT()
{
    no_flow=0
    ,sum_pkts=0,sum_bytes=0,pkt2=0,usrcIP=0,udstIP=0,
    usrcPT=0,udstPT=0;

    r=FR.etime;
    oo=0;
}

```

```

void W_file()//function read from file and write
on other file
{
    fprintf(wf,"%lu",FR.flowId);
    fprintf(wf,"%lu",FR.stime);
    fprintf(wf,"%lu",FR.stmsec);
    fprintf(wf,"%lu",FR.etime);
}

```

```

fprintf(wf, "%lu, ", FR.etmsec);
fprintf(wf, "%u, ", FR.protocol);
fprintf(wf, "%lu, ", FR.srcIP);
fprintf(wf, "%lu, ", FR.srcPT);
fprintf(wf, "%lu, ", FR.dstIP);
fprintf(wf, "%lu, ", FR.dstPT);
fprintf(wf, "%lu, ", FR.pkts);
fprintf(wf, "%lu, ", FR.bytes);
fprintf(wf, "%lu\n", FR.clss);
}

```

```

void parameter()
{

no_flow++;
sum_pkts+= FR.pkts;
sum_bytes+=FR.bytes;
    if(FR.pkts<=2)
        pkt2++;

foundIPadrs = 0;
for (  n=0;n < usrcIP; n++ )
    {
        if(s[n]==FR.srcIP)
        {
            foundIPadrs = 1;

```

```

        break;
    }

}

if (foundIPadrs ==0){
    s[usrcIP]=FR.srcIP;
    usrcIP++;
}

foundIPdst = 0;
for ( a=0;a < udstIP; a++ )
{
    if(y[a]==FR.dstIP)
    {
        foundIPdst = 1;
        break;
    }

}

if (foundIPdst ==0){
    y[udstIP]=FR.dstIP;
    udstIP++;
}

foundptdrs = 0;
for ( b=0;b< usrcPT; b++ )
{
    if(x[b]==FR.srcPT)

```

```

        {
            foundptdrs = 1;
            break;
        }

    }

    if (foundptdrs ==0){
        x[usrcPT]=FR.srcPT;
        usrcPT++;
    }

    foundptadst = 0;
    for ( c=0;c < udstPT; c++ )
    {
        if (v[c]==FR.dstPT)
        {
            foundptadst = 1;
            break;
        }

    }

    if (foundptadst ==0){
        v[udstPT]=FR.dstPT;
        udstPT++;
    }

```

```

if (FR.class!=0)
{
    oo++;
}
}

void writeTAT()
{

    if ((wr=fopen("D:/result.csv",
"a+"))==NULL)//reading file
{
    printf( "writing file could not be opened\n" );
    exit(1);
}
e=(oo*1.0/no_flow) ;

//if(oo>=1)
//    label=1;
//    else
    if(oo>= 50)
        label=1;
    else
        label=0;

fprintf(wr,"%lu,",no_flow );
fprintf(wr,"%lu,",sum_pkts);

```

```
fprintf(wr,"%lu",sum_bytes);  
fprintf(wr,"%lu",usrcIP);  
fprintf(wr,"%lu",udstIP);  
fprintf(wr,"%lu",usrcPT);  
fprintf(wr,"%lu",udstPT);  
fprintf(wr,"%lu",pkt2);  
fprintf(wr,"%lu",oo);  
fprintf(wr,"%lu,\n",label);  
fclose(wr);  
}
```

B.APPENDIX B

RESULT OBTAINED FROM C PROGRAM

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
1	20677	848356	540325054	1408	5519	8579	6379	6973	34	376	1	1	1	0
1	17990	579621	297161470	1547	5656	7752	5693	7376	41	83	1	1	1	0
1	15508	484412	257453804	1374	4657	6715	5044	5657	36	80	1	0	0	0
2	8584	365851	269799515	963	3525	4363	2811	3998	47	219	1	0	0	0
2	17915	507467	286803947	1425	5302	7927	5553	6812	38	165	1	1	1	0
2	17567	625479	380394784	1416	5590	7534	5398	6899	39	159	1	1	1	0
2	18619	644045	501219605	1734	6329	8001	5848	7634	41	189	1	1	1	0
2	20209	651262	406589647	1827	6545	8892	6545	8465	42	82	1	1	1	0
2	19778	1283797	799847061	1623	5842	8753	6335	7275	37	198	1	1	1	0
2	16841	584019	362862599	1515	5773	7333	5093	7263	43	145	1	0	0	0
2	17044	520887	390218228	1496	5745	7167	5232	7367	43	112	1	0	0	0
2	18587	646203	342534678	1448	6101	7886	5473	7675	41	108	1	1	1	0
2	18461	404032	272220206	1439	5822	7825	5401	7284	39	253	1	1	1	0
2	17795	1822400	748844609	1468	5980	8280	4944	7196	40	1128	1	1	1	0
3	10963	1538206	1432941350	1032	4458	6216	3053	5181	47	1027	1	0	0	0
3	18546	492545	311375506	1413	5769	7953	5645	7315	39	266	1	1	1	0
3	19052	944574	719329701	1589	6020	8654	5953	7889	41	694	1	1	1	0
3	19816	550920	361802548	1530	6260	9349	5790	7967	40	1090	1	1	1	0
3	19430	539170	447759029	1463	6005	8647	5572	7416	38	1092	1	1	1	0
3	19714	652063	513076347	1555	5780	8030	5681	7387	37	1102	1	1	1	0
3	18850	421759	288400418	1466	5734	8230	5718	7260	39	562	1	1	1	0
3	19696	431599	282028717	1484	6004	8930	5797	7649	39	679	1	1	1	0

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
3	19589	403088	289407206	1612	6512	8694	5510	8521	43	1212	1	1	1	0
3	17405	478040	366502051	1448	5712	7704	4818	7236	42	1190	1	0	0	0
3	19009	1522670	1227728311	1464	5603	8662	5634	7117	37	929	1	1	1	0
3	17577	711579	448100428	1403	5519	7820	5346	7118	40	154	1	1	1	0
4	19317	475012	356276046	1904	6385	9845	6178	9168	47	2155	1	1	1	0
4	19762	381591	270607009	1857	6361	9635	6477	8574	43	1564	1	1	1	0
4	20089	620034	429908606	1515	6167	9445	6327	9235	46	1848	1	1	1	0
4	19620	531385	352863997	1501	6474	9031	5636	8308	42	965	1	1	1	0
4	19478	1392149	539408241	1485	5896	9813	6152	7786	40	2645	1	1	1	0
4	19159	879486	449489052	1458	6355	9076	5774	8043	42	1557	1	1	1	0
4	19528	525267	335401946	1501	6205	8829	6048	8318	43	930	1	1	1	0
4	21973	551953	419939038	1795	6800	11120	6966	8969	41	1938	1	1	1	0
4	18929	378804	246874300	1623	6586	9302	5968	8665	46	1433	1	1	1	0
4	18058	2666904	1776439584	1399	6059	8482	5586	7905	44	1925	1	0	0	0
4	18941	329687	215168228	1629	6519	9933	6250	8305	44	2405	1	1	1	0
4	18950	359759	224902421	1508	6441	9604	5985	8104	43	1149	1	1	1	0
4	12707	316107	179509984	1150	4402	6422	4218	5332	42	215	1	0	0	0
5	22104	528195	353103371	1310	6082	10582	7063	8567	39	106	1	1	1	0
5	20431	1705664	409907127	1248	6102	9219	6483	8264	40	916	1	1	1	0
5	22291	1073603	627362315	1273	6098	11735	7166	7999	36	2030	1	1	1	0
5	20809	532990	320944204	1337	6461	10272	6537	8349	40	2030	1	1	1	0
5	20307	498737	236897342	1540	6613	9665	6498	8643	43	2153	1	1	1	0

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
5	18672	723199	330145238	1565	6710	8946	5982	8517	46	1555	1	1	1	0
5	19035	443066	297580540	1439	6616	9059	5915	8441	44	1927	1	1	1	0
5	18417	524346	320531844	1279	6304	8782	5505	7773	42	2580	1	0	0	0
5	18570	455112	327703669	1426	6267	8868	5699	8027	43	2476	1	1	1	0
5	1215	28399	16061103	247	731	841	416	623	51	479	1	0	0	0
1	2498	37323	25291825	482	1481	1757	803	1661	66	0	0	0	0	0
1	15945	392211	295215205	1444	5601	7442	4754	7019	44	0	0	0	0	0
1	18257	545033	352867533	1858	6210	8371	5964	8382	46	0	0	0	0	0
1	18407	385685	258533239	1574	5848	8294	5758	7868	43	0	0	0	0	0
1	19341	517467	419437455	1553	6157	8575	5833	8252	43	0	0	1	0	1
1	18367	1514438	1321584451	1652	5967	8175	5722	7816	43	0	0	0	0	0
1	16073	624246	380146654	1706	5773	7453	5141	7436	46	0	0	0	0	0
1	16394	770140	646815980	1488	5514	7597	5244	7061	43	0	0	0	0	0
1	23142	1447094	796137019	1536	5686	9366	7174	7329	32	0	0	1	0	1
1	24286	1233385	328449049	1544	6072	9705	7484	7977	33	0	0	1	0	1
1	22297	1380669	877136387	1482	5455	9221	7101	7130	32	0	0	1	0	1
1	818	34720	10888481	239	464	535	353	369	45	0	0	0	0	0
2	16530	504356	351770353	1331	5374	7157	4973	6885	42	19	0	0	0	0
2	16453	1342483	579417990	1400	5425	7083	4875	6978	42	2	0	0	0	0
2	17423	1813337	473438841	1492	5603	7382	5213	7098	41	0	0	0	0	0
2	8194	1220715	1183308214	894	1762	3957	3376	1668	20	4	0	0	0	0

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
3	18358	396417	280491027	1420	5824	7509	5471	7527	41	0	0	0	0	0
3	11866	283959	159097466	1207	3999	5244	3829	4773	40	0	0	0	0	0
3	129	1707	558736	61	72	67	80	24	19	0	0	0	0	0
4	3428	81102	59611811	605	1827	2142	1121	2019	59	0	0	0	0	0
4	61	1563	1081848	31	34	37	31	7	11	0	0	0	0	0
5	13158	264790	126838072	1043	4754	7935	4087	5985	45	3	0	0	0	0
5	17297	346091	194685648	1336	6341	8962	5126	8086	47	3	0	0	0	0
5	18846	1511556	1251900453	1324	6265	9942	5668	8549	45	1	0	0	0	0
5	18848	2454892	786913931	1371	6370	9535	5699	8122	43	4	0	0	0	0
6	11088	203179	118178983	933	4179	5790	3541	4807	43	3	0	0	0	0
6	19106	404486	241163465	1276	6374	8392	5591	8039	42	3	0	0	0	0
6	18893	588556	438847925	1211	6266	8775	5599	7746	41	7	0	1	0	1
6	19382	471304	321743722	1351	6499	8769	5798	8090	42	0	0	0	0	0
6	17836	578234	271201400	1218	6412	7759	5164	7983	45	4	0	0	0	0
6	18619	518570	347672056	1300	6523	8590	5382	8083	43	1	0	0	0	0
6	19626	893806	632052060	1620	6907	9452	6336	8921	45	4	0	1	0	1
6	18057	391309	286653967	1374	6171	8383	5734	7774	43	0	0	0	0	0
6	17914	473343	290946817	1265	6063	7898	5493	7545	42	3	0	0	0	0
6	16714	787490	633110569	1203	6013	7861	5051	7333	44	2	0	0	0	0
6	16829	357956	228596400	1207	6066	8104	4964	7473	44	0	0	0	0	0
6	17851	466487	262707460	1598	6758	8477	5557	8327	47	3	0	0	0	0
6	15970	613983	443316073	1299	6253	7441	4718	7506	47	0	0	0	0	0

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
6	16583	392218	235493516	1326	6460	7754	5135	8077	49	0	0	0	0	0
6	5532	139935	88447672	699	1845	3187	2331	1807	33	1	0	0	0	0
7	1505	81948	20416801	371	879	1114	501	1083	72	2	0	0	0	0
7	13286	479297	350884329	1258	6224	6615	3725	7541	57	0	0	0	0	0
7	14986	509368	324742978	1391	6284	6722	4152	7890	53	2	0	0	0	0
7	16650	1096622	578337379	1317	6619	7345	4716	8182	49	1	0	0	0	0
7	15532	572566	498489831	1159	6220	7181	4331	7552	49	0	0	0	0	0
7	15541	331208	256741710	1245	6508	7198	4314	8008	52	4	0	0	0	0
7	14854	520562	402675349	1227	6782	6533	3848	8139	55	0	0	0	0	0
7	14795	458990	344741307	1154	6709	6608	3801	7882	53	2	0	0	0	0
7	16890	453502	335407114	1236	6605	6997	4667	8100	48	1	0	0	0	0
7	17429	407840	309879703	1227	6836	7298	4749	8528	49	0	0	0	0	0
7	17333	816504	691355604	1324	7140	7220	4596	9045	52	2	0	0	0	0
7	15725	397191	259950558	1223	7002	7301	4027	8352	53	0	0	0	0	0
7	15864	421433	312232234	1263	6799	6465	4054	8335	53	4	0	0	0	0
7	16337	404915	264546592	1181	6673	6868	4348	8181	50	0	0	0	0	0
7	16650	438777	266068118	1314	6988	7583	4729	8489	51	1	0	0	0	0
7	17161	673578	366135116	1511	7475	8004	4788	9663	56	2	0	0	0	0
7	9462	263510	214345408	979	3515	4758	3172	3903	41	0	0	0	0	0
8	12922	287026	220686518	1108	6459	5522	3296	7568	59	2	0	0	0	0
8	16139	520941	430371911	1177	6883	6837	3996	8519	53	1	0	0	0	0

TAT_ID	flows	pckts	bytes	srcIP	dstIP	srcPT	dstPT	2pflow	2pflow/flows	mal	label	RULE	TP	FP
8	17494	787629	653047918	1254	6912	7966	4680	8308	47	0	0	0	0	0
8	16306	413223	292545094	1227	7012	6951	4114	8442	52	0	0	0	0	0
8	15934	1051463	1089584825	1281	6783	6585	4209	8331	52	0	0	0	0	0
8	16473	642484	286619802	1205	6938	7258	4308	8312	50	0	0	0	0	0
8	15775	589599	300784703	1133	6621	7282	4007	8239	52	0	0	0	0	0
8	16695	602730	278897830	1162	6977	7187	4309	8616	52	0	0	0	0	0
8	17155	1075163	1074631170	1144	6883	7213	4558	8521	50	0	0	0	0	0
8	16772	1560959	643634764	1205	6932	7662	4496	8448	50	0	0	0	0	0
8	16452	505609	428453835	1151	6561	7378	4251	8363	51	0	0	0	0	0
8	16852	845286	780282119	1252	6755	7261	4555	8273	49	0	0	0	0	0
8	12871	673403	521157259	1046	5192	6425	3732	5771	45	0	0	0	0	0
8	12149	763155	701345151	994	5188	6490	3528	5909	49	0	0	0	0	0
8	16650	562686	384317786	1215	6972	7497	4301	8577	52	0	0	0	0	0
8	13271	601277	380246401	1073	4370	6391	4274	5189	39	0	0	0	0	0
8	90	3462	1034443	39	81	72	36	64	71	0	0	0	0	0
											49		45	39
avg	16260.16	656712.5	421529215	1293.74	5673.496	7465.561	4858.024	7135.772	44.36585366				0.8	0.08
std	4929.49	461359.6	297193341	346.3115	1599.467	2212.683	1533.568	2111.073	8.076750983					